

Math 303, Fall 2011, Lecture 11

① The symbols of a formal language (Cohen I.2 (starting on p3))

Cohen talks of two kinds of symbols

general symbols common to all mathematics  
and

symbols which are special to the kind of math  
we're doing

If we wish to use set theory as a foundation for all  
of mathematics then we only need special symbols for set  
theory. Every other special symbol will just be an  
abbreviation for something which can be said  
using set theory, albeit perhaps in a very convoluted way.

So, our only special symbol is

$\in$  element of

And our general symbols are

$\sim$  (sometimes  $\neg$ ) not

$\wedge$  (sometimes  $\&$ ) and

$\vee$  or

$\rightarrow$  implies

$\leftrightarrow$  if and only if

$\forall$  for all

$\exists$  there exists

$(, )$  parentheses  
 $=$  equality

$x_0, x_1, x_2, \dots, x, y, z, w, \dots$

$c_0, c_1, \dots$

Note you can remember this  
and because it looks like  
an A

Note elsewhere in math  
usually use double  
arrows  $\Rightarrow, \Leftrightarrow$   
in logic always single  
arrows

variable symbols

constant symbols

# Notes

①

Cohen also has names for relations  
but we don't need this because we can form  
relations as sets (just as for functions)

②

We gave English intuitive meanings for each symbol  
but we'll give rules independent of the meanings  
to manipulate these. But keep the intuitive meanings in  
mind to have intuition of how to manipulate  
formally

③

If you want your language to be finite  
we need countably many variable symbols  
from finitely many actual symbols

solution variables  $x, x', x'', x''', \dots$

same for constants  $c, c', c'', c''', \dots$

eg  $x \subseteq y$

$$\forall x' ((x' \in x) \rightarrow (x' \in y))$$

eg lets allow ourselves + in its usual sense for natural numbers  
(could define it as a function using sets)

$$x + y = y + x \text{ for all } x, y \in \omega$$

$$\forall x \forall y ((x \in \omega \wedge y \in \omega)$$

$$\rightarrow (x + y = y + x))$$

here I'm using  $\omega$  as a constant symbol

eg The empty set

$$\exists x (\forall x' (\neg(x' \in x)))$$

One of our axioms tells us such an  $x$  exists  
so this statement is true

Knowing this - we can avoid having to say all this  
anytime we need  $\emptyset$

Just use  $\phi$  knowing I could put it in the whole statement  
This is what is meant by an abbreviation

They are part of our formal system but they  
make things easier to read.

## ② The syntax of the formal system

Some things I can write with these symbols make sense

$$\forall z (z \in x \rightarrow z \in y)$$

Others do not

$$\forall \rightarrow) x x x x z \in (( \rightarrow$$

But we want to be able to use these symbols  
without appealing to their meaning. So we need  
rules for forming valid strings

The valid strings are called well formed formulas (wff)

Here are the rules

① For any  $x, y$  variable symbols and any  $c, d$  constant symbols. The following are well formed formulas

$(x \in y)$

$(x \in c)$

$(c \in x)$

$(c \in d)$

Note: These aren't necessarily true, just syntactically ok  
eg  $(x \in \emptyset)$

② For  $x, y, c, d$  as above. The following are well formed formulas

$(x = y)$        $(c = d)$

$(x = c)$

$(c = x)$

③ If  $A$  and  $B$  are well formed formulas then so are

$$(\neg A) \quad (A \wedge B) \quad (A \vee B) \quad (A \rightarrow B) \\ (A \leftrightarrow B)$$

④ If  $A$  is a well formed formula and  $x$  is a variable symbol then

$$(\forall x A) \quad (\exists x A)$$

Note again well formed formulas don't need to be true

eg  $(\forall x (\forall y (x=y)))$

is a well formed formula

But not true in set theory

to see this is well formed. ①  $(x=y)$  is well formed  
so ②  $(\forall y (x=y))$  is well formed so ③  $(\forall x (\forall y (x=y)))$  is well formed

Note Well leave out parentheses when the meaning is clear

eg  $\forall x \forall y (x=y)$  is clear

$\forall x \forall y x=y$  is not clear enough

eg Are the following well formed? (don't worry about parentheses if it's clear)

$x=y$       yes      good

$x=y=z$       no       $(x=y) \wedge (y=z)$       yes a wff

$\forall x (x \wedge y)$        $\forall x A$  where  $A = x \wedge y$ , but  $x \wedge y$  is not a well formed formula

$\forall x \exists y \exists z (y=z)$       yes      not good      var  $x$  is not a free var in  $\exists y \exists z (y=z)$

③ Next time

- free and bound variables, sentences
- examples
- truth

Read Cohen Ch I sections 2 and 3