

Math 343, lecture 13

## ① Practicalities of Boltzmann samplers

Some important questions regarding Boltzman samplers

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## ② Speed

As a first attempt to answer the speed question  
lets forget about

Each piece

provided

oracle assumption

It is plausible in practice as

This brings up another question

.

The above discussion gives

Proposition

let  $A$  be a combinatorial class with a specification (either iterative or recursive) in terms of the constructions we know (and some others also work)

Then under the oracle assumption the generation of a  $\in A$  by the Boltzmann generate of  $A$  takes

What about the oracle itself and sensitivity to round off errors.

### ③ Geometric random numbers

How to implement a geometric random number generator assuming you have a uniform random number generator:

Algorithm

Geometric-rand

input  $\lambda$  (parameter to geometric distribution)

$$p(k) = (1-\lambda)\lambda^k$$

$v = \text{rand}()$

(uniform random in  $(0, 1)$ )

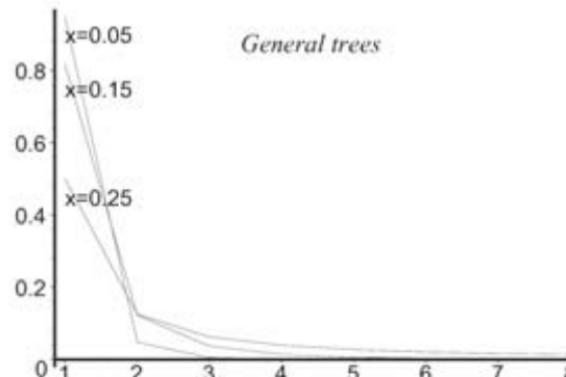
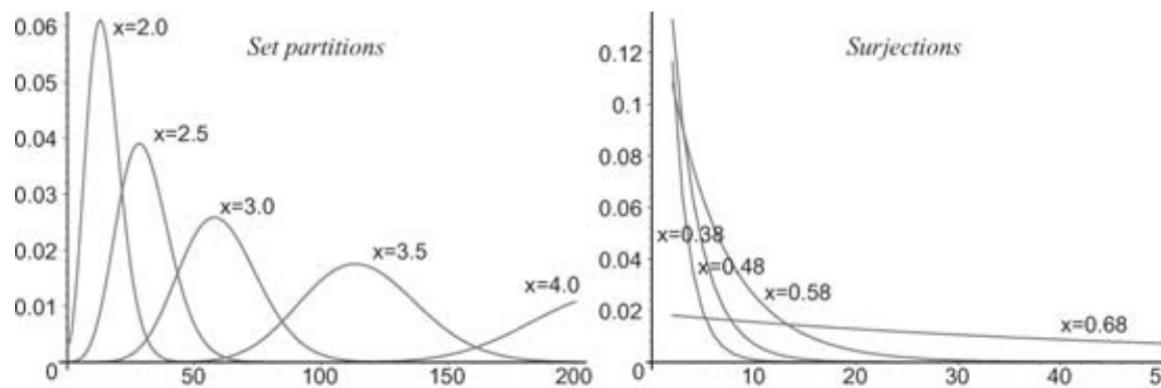
This approach will also work for other distributions just change  $p(k)$ .

#### ④ Distributions and finding the optimal $x$

As we saw last time the expected size of the Boltzmann model with parameter  $x$  for a class  $\mathcal{C}$  is

$$E_x = x \frac{A'(x)}{A(x)} \quad \text{so}$$

But



The bumpy distributions are the best for Boltzmann sampling and there are precise conditions which guarantee that they occur.

Typical examples are

What about flat distributions → we can use a rejection strategy

With a flat distribution this strategy will succeed

Classes of words typically have flat distributions.

Finally the most troublesome ones — peaked distributions

## ⑤ Pointing

Consider a combinatorial class  $\mathcal{C}$  built recursively out of  $+, \times, \mathcal{E}, \mathcal{Z}$ , Seq

Build a new class  $\mathcal{C}^*$

Each  $c \in \mathcal{C}$  gives

so

$$\mathcal{C}^*(x) =$$

pointing

Using the rules for derivatives we can translate our specifications to the pointed classes

$\mathcal{Z}^*$

$\mathcal{E}^*$

$$(A + B)^* =$$

$$(A \times B)^* =$$

$$(\text{Seq } A)^* =$$

eg

for the Boltzmann generator

Let's see it

⑥ Next time

Gray codes